## UCF

## Stands For Opportunity

CDA6530: Performance Models of Computers and Networks

## Chapter 5: Generating Random Number and Random Variables

## Objective

- Use computers to simulate stochastic processes
- Learn how to generate random variables - Discrete r.v.
- Continuous r.v.
- Basis for many system simulations


## Pseudo Random Number Generation (PRNG)

a $x_{n}=a x_{n-1} \bmod m$

- Multiplicative congruential generator
- $X_{n}=\{0,1, \cdots, m-1\}$
- $x_{n} / m$ is used to approx. distr. $U(0,1)$
a $x_{0}$ is the initial "seed"
- Requirements:
- No. of variables that can be generated before repetition begins is large
- For any seed, the resultant sequence has the "appearance" of being independent
- The values can be computed efficiently on a computer
- $x_{n}=a x_{n-1} \bmod m$
a $m$ should be a large prime number
- For a 32-bit machine (1 bit is sign)
- $m=2^{31}-1=2,147,483,647$
- $a=7^{5}=16,807$
- For a 36-bit machine
a $m=2^{35}-31$
- $a=5^{5}$
a $x_{n}=\left(a x_{n-1}+c\right) \bmod m$
$\square$ Mixed congruential generator


## In C Programming Language

- Int rand(void)
- Return int value between 0 and RAND_MAX - RAND_MAX default value may vary between implementations but it is granted to be at least 32767
- $X=$ rand ()
- $\mathrm{X}=\{0,1, \cdots$, RAND_MAX $\}$
- $X=$ rand ()$\% m+n$
- $\mathrm{X}=\{\mathrm{n}, \mathrm{n}+1, \cdots, \mathrm{~m}+\mathrm{n}-1\}$
- Suitable for small m;
- Lower numbers are more likely picked


## (0,1) Uniform Distribution

- $\mathrm{U}(0,1)$ is the basis for random variable generation
- C code (at least what I use): Double rand01) \{ double temp; temp $=$ double $($ rand ()$+0.5) /$ (double(RAND_MAX) + 1.0); return temp;


## Generate Discrete Random Variables ---- Inverse Transform Method

a r.v. $X: P\left(X=x_{j}\right)=p_{j}, \quad j=0,1, \cdots$
$\square$ We generate a PRNG value $\mathrm{U} \sim \mathrm{U}(0,1)$

- For $0<\mathrm{a}<\mathrm{b}<1, \mathrm{P}(\mathrm{a} \leq \mathrm{U}<\mathrm{b}\}=\mathrm{b}-\mathrm{a}$, thus

$$
P\left(X=x_{j}\right)=P\left(\sum_{i=0}^{j-1} p_{i} \leq U<\sum_{i=0}^{j} p_{i}\right)=p_{j}
$$

$$
X= \begin{cases}x_{0} & \text { if } U<p_{0} \\ x_{1} & \text { if } p_{0} \leq U<p_{0}+p_{1} \\ \vdots & \\ x_{j} & \text { if } \sum_{i=0}^{j-1} p_{i} \leq U<\sum_{i=0}^{j} p_{i}\end{cases}
$$

## Example

- A loaded dice: $\square P(1)=0.1 ; P(2)=0.1 ; P(3)=0.15 ; P(4)=0.15$ - $P(5)=0.2 ; ~ P(6)=0.3$
- Generate 1000 samples of the above loaded dice throwing results


## Generate a Poisson Random Variable

$$
p_{i}=P(X=i)=e^{-\lambda \frac{\lambda^{i}}{i!}}, \quad i=0,1, \cdots
$$

- Use following recursive formula to save computation:

$$
p_{i+1}=\frac{\lambda}{i+1} p_{i}
$$

## Some Other Approaches

- Acceptance-Rejection approach
- Composition approach
- They all assume we have already generated a random variable first (not U)
$\square$ Not very useful considering our simulation purpose


## Generate Continuous Random Variables ---- Inverse Transform Method

a r.v. $X: F(x)=P(X \leq x)$
a r.v. $Y: Y=F^{-1}(U)$

- $Y$ has distribution of $F$. $\left(Y={ }_{\text {st }} X\right)$
- $\mathrm{P}(\mathrm{Y} \leq \mathrm{x})=\mathrm{P}\left(\mathrm{F}^{-1}(\mathrm{U}) \leq \mathrm{x}\right)$
$=P\left(F\left(F^{-1}(U)\right) \leq F(x)\right)$
$=P(U \leq F(x))$
$=P(X \leq x)$
- Why? Because $0<F(x)<1$ and the CDF of a uniform $\mathrm{F}_{\mathrm{U}}(\mathrm{y})=\mathrm{y}$ for all $\mathrm{y} \in[0 ; 1]$


## Generate Exponential Random Variable

$$
\begin{gathered}
F(x)=1-e^{-\lambda x} \\
U=1-e^{-\lambda x} \\
e^{-\lambda x}=1-U \\
x=-\ln (1-U) / \lambda \\
F^{-1}(U)=-\ln (1-U) / \lambda
\end{gathered}
$$

## Generate Normal Random Variable --- Polar method

- The theory is complicated, we only list the algorithm here:
- Objective: Generate a pair of independent standard normal r.v. ~ N(0, 1)
- Step 1: Generate $(0,1)$ random number $U_{1}$ and $U_{2}$
- Step 2: Set $\mathrm{V}_{1}=2 \mathrm{U}_{1}-1, \mathrm{~V}_{2}=2 \mathrm{U}_{2}-1 \quad S=V_{1}^{2}+V_{2}^{2}$
- Step 3: If S> 1, return to Step 1.
- Step 4: Return two standard normal r.v.:

$$
X=\sqrt{\frac{-2 \ln S}{S}} V_{1}, \quad Y=\sqrt{\frac{-2 \ln S}{S}} V_{2}
$$

| $\mathbf{z}$ | $F(x)$ | $\mathbf{z}$ | $F(X)$ | $\mathbf{z}$ | $F(X)$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| -2.5 | 0.006 | -1 | 0.159 | 0.5 | 0.691 |
| -2.4 | 0.008 | -0.9 | 0.184 | 0.6 | 0.726 |
| -2.3 | 0.011 | -0.8 | 0.212 | 0.7 | 0.758 |
| -2.2 | 0.014 | -0.7 | 0.242 | 0.8 | 0.788 |
| -2.1 | 0.018 | -0.6 | 0.274 | 0.9 | 0.816 |
| -2 | 0.023 | -0.5 | 0.309 | 1 | 0.841 |
| -1.9 | 0.029 | -0.4 | 0.345 | 1.1 | 0.864 |
| -1.8 | 0.036 | -0.3 | 0.382 | 1.2 | 0.885 |
| -1.7 | 0.045 | -0.2 | 0.421 | 1.3 | 0.903 |
| -1.6 | 0.055 | -0.1 | 0.46 | 1.4 | 0.919 |
| -1.5 | 0.067 | 0 | 0.5 | 1.5 | 0.933 |
| -1.4 | 0.081 | 0.1 | 0.54 | 1.6 | 0.945 |
| -1.3 | 0.097 | 0.2 | 0.579 | 1.7 | 0.955 |
| -1.2 | 0.115 | 0.3 | 0.618 | 1.8 | 0.964 |
| -1.1 | 0.136 | 0.4 | 0.655 | 1.9 | 0.971 |

- Another approximate method- Table lookup
- Treat Normal distr. r.v. X as discrete r.v.
- Generate a $U$, check $U$ with $F(x)$ in table, get $z$


## Generate Normal Random Variable

- Polar method generates a pair of standard normal r.v.s X~N(0,1)
- What about generating r.v. $\mathrm{Y} \sim \mathrm{N}\left(\mu, \sigma^{2}\right)$ ?

$$
\square \mathrm{Y}=\sigma \mathrm{X}+\mu
$$

## Generating a Random Permutation

- Generate a permutation of $\{1, \cdots, n\}$
- $\operatorname{Int}(\mathrm{kU})+1$ :
- uniformly pick from $\{1,2, \cdots, k\}$
- Algorithm:
- $P_{1}, P_{2}, \cdots, P_{n}$ is a permutation of $1,2, \cdots, n$ (e.g., we can let $P_{j}=j, j=1, \cdots, n$ )
- Set $\mathrm{k}=\mathrm{n}$
- Generate U, let I $=\operatorname{lnt}(\mathrm{kU})+1$
- Interchange the value of $P_{1}$ and $P_{k}$
- Let $\mathrm{k}=\mathrm{k}-1$ and if $\mathrm{k}>1$ goto Step 3
$\square P_{1}, P_{2}, \cdots, P_{n}$ is a generated random permutation
Example: permute (10, 20, 30, 40, 50)


## Monte Carlo Approach ----

## Use Random Number to Evaluate Integral

$\theta=\int_{0}^{1} g(x) d x$

$$
\theta=E[g(U)]
$$

$\square U$ is uniform distr. r.v. $(0,1)$ - Why?
$E[X]=\int_{-\infty}^{\infty} x f(x) d x$

$$
\begin{gathered}
E[g(X)]=\int_{-\infty}^{\infty} g(x) f(x) d x \\
f_{U}(x)=1 \quad \text { if } 0<x<1
\end{gathered}
$$

- $\mathrm{U}_{1}, \mathrm{U}_{2}, \cdots, \mathrm{U}_{\mathrm{k}}$ are independent generated uniform distr. $(0,1)$
- $g\left(U_{1}\right), \cdots, g\left(U_{k}\right)$ are independent
- Law of large number:

$$
\sum_{i=1}^{k} \frac{g\left(U_{i}\right)}{k} \rightarrow E[g(U)]=\theta \text { as } k \rightarrow \infty
$$

$$
\theta=\int_{a}^{b} g(x) d x
$$

- Substitution: $y=(x-a) /(b-a), d y=d x /(b-a)$

$$
\begin{gathered}
\theta=\int_{0}^{1}(b-a) \cdot g(a+(b-a) y) d y=\int_{0}^{1} h(y) d y \\
h(y)=(b-a) \cdot g(a+(b-a) y)
\end{gathered}
$$

$$
\theta=\int_{0}^{1} \int_{0}^{1} \cdots \int_{0}^{1} g\left(x_{1}, \cdots, x_{n}\right) d x_{1} d x_{2} \cdots d x_{n}
$$

$$
\theta=E\left[g\left(U_{1}, \cdots, U_{n}\right)\right]
$$

- Generate many g(....) - Compute average value
- which is equal to $\theta$

